ROCKSET: The design and implementation of a data system for low-latency queries for search and analytics

**Dhruba Borthakur, Igor Canadi** 

**Rockset** 



## **Speakers**



#### **Dhruba**

- CTO and Co-Founder at Rockset
- RocksDB at Facebook
- Hadoop File System
- HBase



#### **Igor**

- Software Engineer at Rockset
- RocksDB at Facebook
- GraphQL



#### **Overview**

- 1. Converged indexing
- 2. Query execution over a distributed index
- 3. High-throughput index updates
- 4. Scaling in the cloud



#### **Motivation**

#### Ease of use

- Minimize configuration
- Can connect to any data source
- Real time writes

#### **High performance**

- Low latency queries
- High throughput writes



# **Converged indexing**



#### **Columnar storage**

- Store each column separately
- **Great compression**
- Only fetch columns query needs









#### **Columnar storage**

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- Great compression
- Only fetch columns query needs

```
<doc 0>
  "name": "Igor",
 "interests": ["databases", "snowboarding"],
  "last active": 2019/3/15
<doc 1>
  "name": "Dhruba",
  "interests": ["cars", "databases"],
  "last active": 2019/3/22
```

#### "name"

0	Igor
1	Dhruba

#### "interests"

0.0	databases
0.1	snowboarding
1.0	cars
1.1	databases

#### "last\_active"

0	2019/3/15
1	2019/3/22



#### **Columnar storage**

#### **Advantages**

- Cost effective
- Narrow queries, wide tables
- Scan queries
- Analytical queries

#### **Disadvantages**

- High write latency
- High minimum read latency
- Not suitable for online applications



### **Search indexing**

- For each value, store documents containing that value (posting list)
- Quickly retrieve a list of document IDs that match a predicate









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- Quickly retrieve a list of document IDs that match a predicate

```
"name"
<doc 0>
                                                                Dhruba 1
 "name": "Igor",
                                                                lgor
                                                                       0
 "interests": ["databases", "snowboarding"],
 "last active": 2019/3/15
                                                                "interests"
                                                                                  0.0; 1.1
                                                                databases
<doc 1>
                                                                                  1.0
                                                                cars
                                                                snowboarding
                                                                                  0.1
 "name": "Dhruba",
 "interests": ["cars", "databases"],
                                                                "last active"
 "last active": 2019/3/22
                                                                2019/3/15
                                                                                  0
                                                               2019/3/22
```



#### **Search indexing**

#### **Advantages**

- High selectivity queries
- Low latency queries
- Suitable for online applications

#### **Disadvantages**

• Slower analytical queries



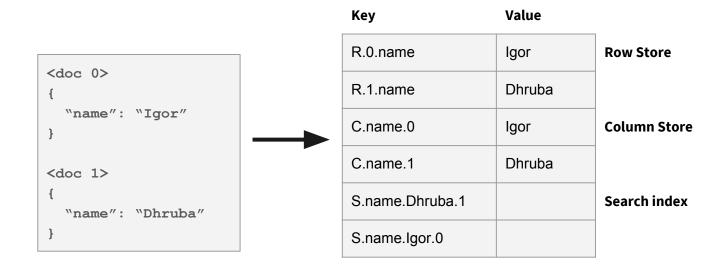
### **Converged indexing**

- Columnar and search indexes in the same system
- Built on top of key-value store abstraction
- Each document maps to many key-value pairs



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## **Converged indexing - queries**

- Fast analytical queries + fast search queries
- Optimizer picks between columnar store or search index



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- Fast analytical queries + fast search queries
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```
SELECT *
FROM search_logs
WHERE keyword = 'strata'
AND locale = 'en'
```

Search index

SELECT keyword, count(\*)
FROM search\_logs
GROUP BY keyword
ORDER BY count(\*) DESC

Columnar store



### Rockset - data system built on converged indexing

- All fields indexed
- SQL
- Document model, schemaless
- Real-time writes, updates and deletes
- Cloud service

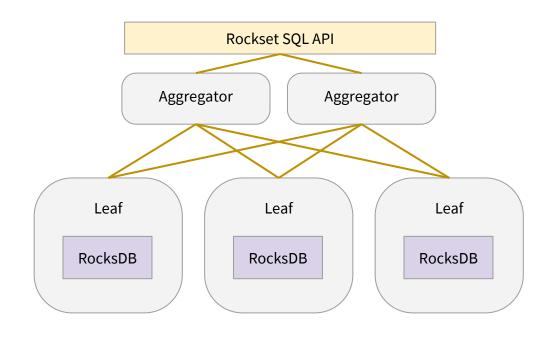




# Query execution over a distributed index

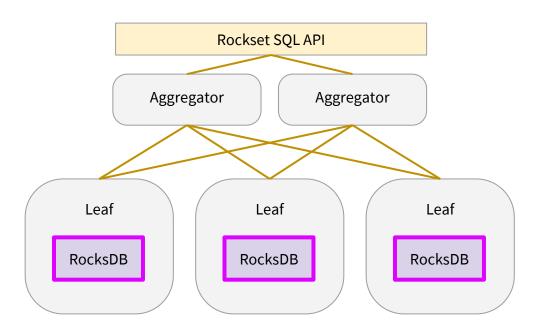


- API layer
- Aggregators
- Leaves



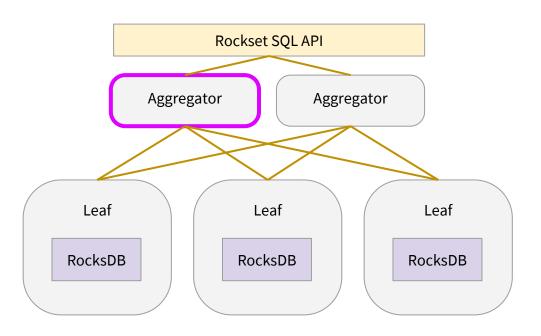


- Converged index sharded across RocksDB instances
- Document-based sharding
- Query hits all shards



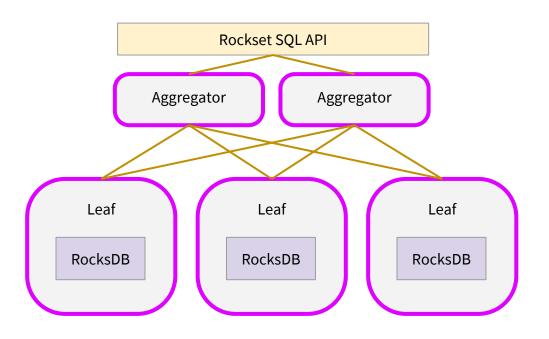


- Aggregator receives the query
- Parsing, compilation, optimization
- Produces operator DAG



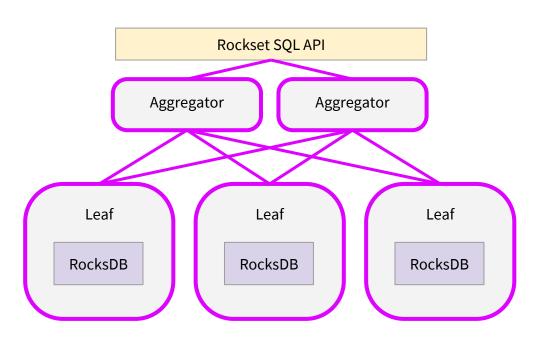


- Set of instructions for each node:
  - Predecessors
  - Operators
  - Successors





- Set of instructions for each node:
  - Predecessors
  - Operators
  - Successors
- On each node:
  - Wait for predecessors
  - Connect to successors
  - Data starts flowing bottom-up





# High-throughput index updates



### Challenges with updating multiple indexes

Maintaining multiple indexes adversely impacts write throughput



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  - Requires consensus coordination between servers



#### Challenges with updating multiple indexes

- Maintaining multiple indexes adversely impacts write throughput
- Challenge 1: one new record = multiple servers updates
  - Requires consensus coordination between servers
- Challenge 2: one new field = multiple random writes
  - Requires increased disk I/O



#### Challenge 1: one new record = multiple servers updates

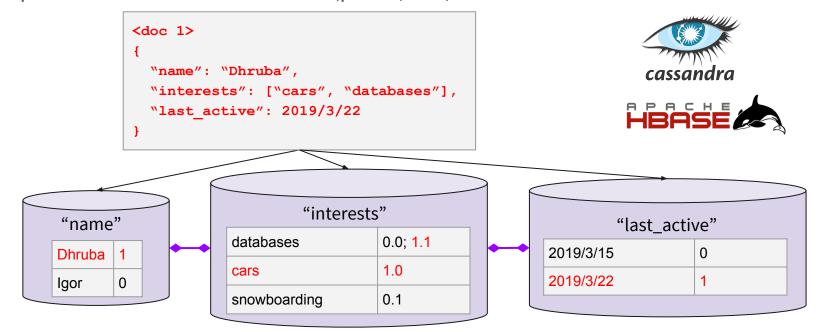
 In a traditional database with term sharding and n indexes, one write incurs updates to n different indexes on n servers

```
<doc 1>
            "name": "Dhruba",
            "interests": ["cars", "databases"],
            "last active": 2019/3/22
                               "interests"
"name"
                                                                        "last active"
                                          0.0; 1.1
                      databases
                                                               2019/3/15
                                                                                   0
Dhruba
                                          1.0
                      cars
                                                               2019/3/22
Igor
                      snowboarding
                                          0.1
```



#### Challenge 1: one new record = multiple servers updates

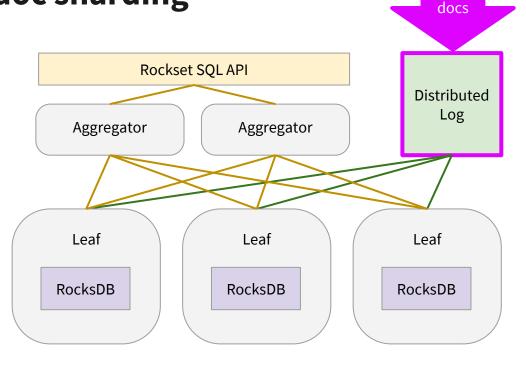
- In a traditional database with term sharding and n indexes, one write incurs updates to n different indexes on n servers
- Requires a distributed transaction (paxos, raft) between **n** servers





# Addressing challenge 1: doc sharding

 Updates are durably-buffered to a distributed log

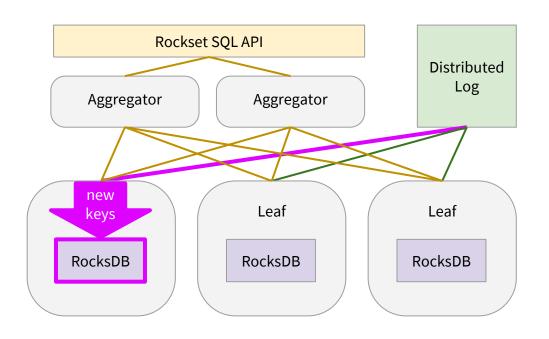




new

## Addressing challenge 1: doc sharding

- Updates are durably buffered to a distributed log
- Leafs tail only documents in the shards they are responsible for
- Doc sharding means all new keys will only affect a single shard/leaf





## Challenge 2: one new doc = multiple random writes

 One doc becomes multiple keys in the index

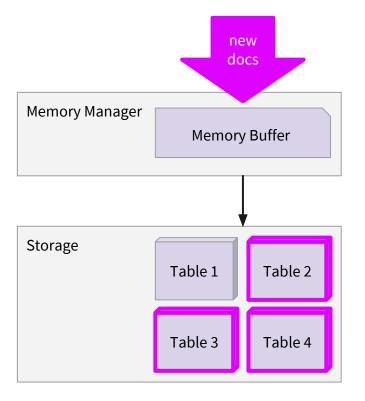
Key	Value	
R.0.name	Igor	Row Store
R.1.name	Dhruba	
C.name.0	Igor	Column Store
C.name.1	Dhruba	
S.name.Dhruba.1		Search index
S.name.lgor.0		





### Challenge 2: one new doc = multiple random writes

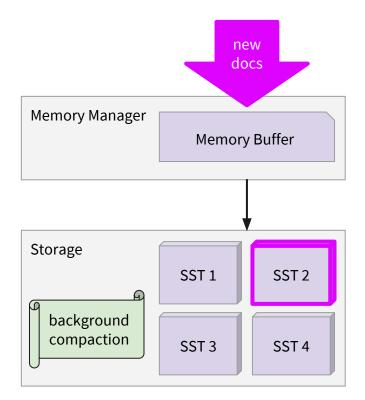
- Traditional systems use B-tree storage structure
- Keys are sorted across tables
- A single record update would incur writes to multiple different tables





### Addressing challenge 2: RocksDB LSM

- RocksDB uses log-structured merge-tree (LSM)
- Multiple record updates accumulate in memory and written into a single SST file
- Keys are sorted between SST files via compaction in a background process
- Multiple index updates from multiple docs result in one write to storage





# Scaling in the cloud



# Key insight into economics of cloud

Cost of 1 cpu for 100 minutes == Cost of 100 cpu for 1 minute!!



### Key insight into economics of cloud

- Cost of 1 cpu for 100 minutes == Cost of 100 cpu for 1 minute!!
  - Without cloud: statically provision for peak demand
  - With cloud: dynamically provision for current demand



#### Key insight into economics of cloud

- Cost of 1 cpu for 100 minutes == Cost of 100 cpu for 1 minute!!
  - Without cloud: statically provision for peak demand
  - With cloud: dynamically provision for current demand
- Goal: scale up and down storage as needed to achieve desired performance



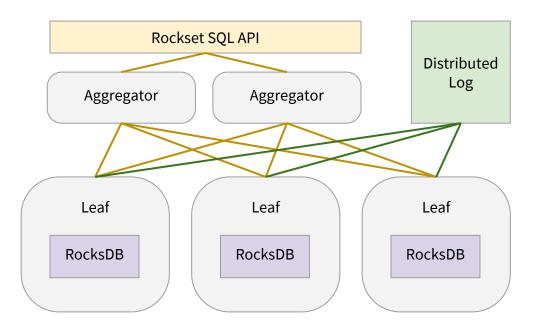
#### **Scheduling compute**

- Kubernetes Horizontal Pod Autoscaler (HPA) to schedule compute pods across pool of nodes (AWS EC2 machines)
  - Using cpu and memory thresholds
- Custom scheduler to manage underlying nodes
  - Manage both demand (of workload) and supply (of hardware)
  - Sum CPU usage of all pods to spin up new nodes when needed
  - Aggressively shed nodes when not needed



# **Scheduling storage**

Each leaf running RocksDB stores indices





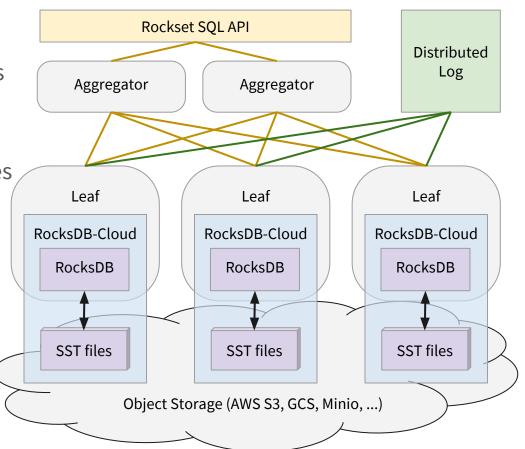
# **Scheduling storage**

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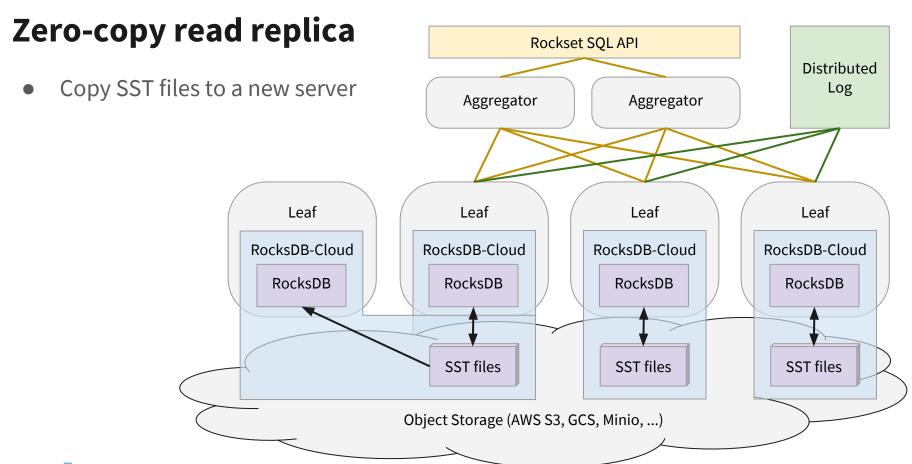
RocksDB-Cloud (open source)
 extends RocksDB to flush SST files
 to durable cloud storage

 RocksDB-Cloud = serverless storage

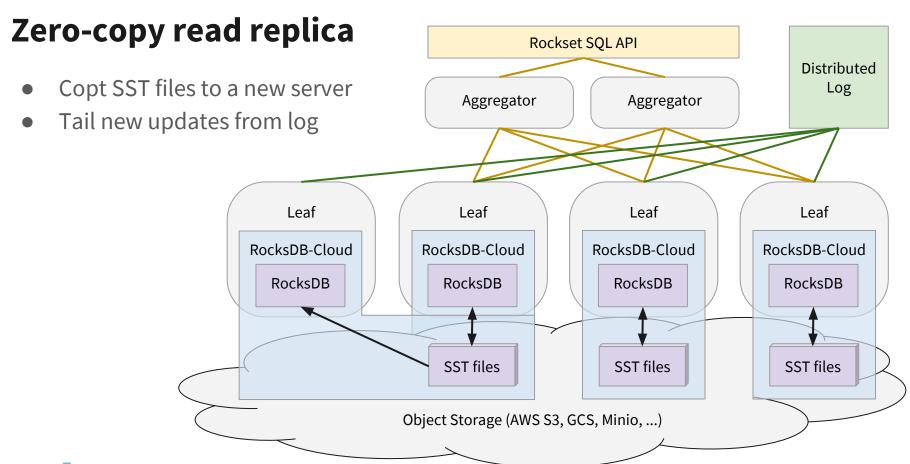
http://github.com/rockset/rocksdb-cloud



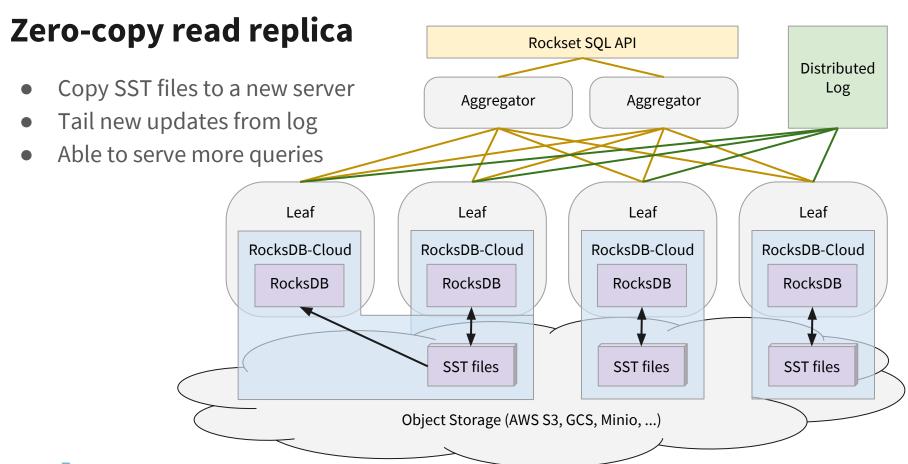






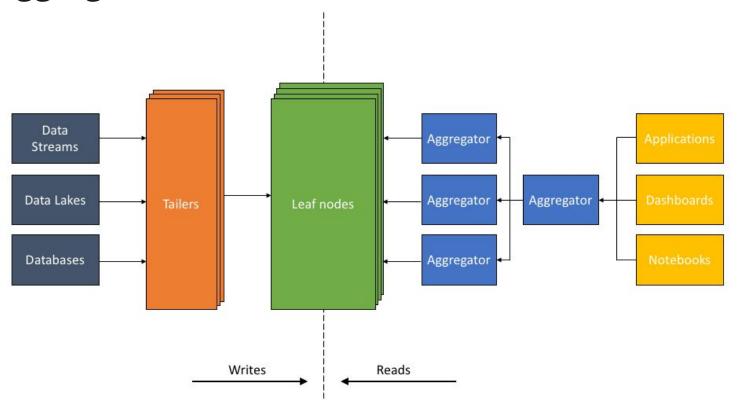






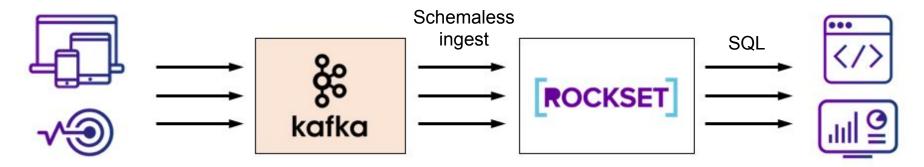


#### The Aggregator Leaf Tailer Architecture





### **Event Analytics Using Rockset**



Events from devices and sensors Applications and live dashboards

- All fields and values indexed
- Low-latency queries
- Serving layer for online applications and live dashboards



# **Check it out: rockset.com**

dhruba@rockset.com igor@rockset.com

Thank you.





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